Naruki Masuda, Hugh Thomas, Andy Tonks, & Bruno Vallette

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THE DIAGONAL OF THE ASSOCIAHEDRA

by Naruki Masuda, Hugh Thomas, Andy Tonks & Bruno Vallette

Abstract. — This paper introduces the first general method to solve the problem of the approximation of the diagonal for face-coherent families of polytopes. We recover the classical cases of the simplices and the cubes and we solve it for the associahedra, also known as Stasheff polytopes. We show that it satisfies an easy-to-state cellular formula. For the first time, we endow a family of realizations of the associahedra (the Loday realizations) with a topological and cellular operad structure; it is shown to be compatible with the diagonal maps.

Résumé (La diagonale de l’associaèdre). — Cet article introduit pour la première fois une méthode générale permettant de résoudre le problème de l’approximation de la diagonale de familles de polytopes satisfaisant à une propriété de cohérence par faces. On retrouve les cas classiques des simplexes et des cubes et on résout celui des associaèdres, appelés aussi polytopes de Stasheff. On montre que ce dernier cas vérifie une formule cellulaire facile à énoncer. Pour la première fois, nous munissons une famille de réalisations des associaèdres (celle de Loday) d’une structure d’opérade topologique cellulaire, dont nous montrons qu’elle est compatible avec les diagonales.

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Introduction

The present paper has three goals: to introduce, for the first time, a general machinery to solve the problem of the approximation of the diagonal of face-coherent families of polytopes (Section 2), to give a complete proof for the case of the associahedra (Theorem 1) and, last but not least, to popularize the resulting magical formula (Theorem 2) to facilitate its application in other domains.

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The problem of the approximation of the diagonal of the associahedra lies at the crossroads of three clusters of domains. There are first the mathematicians who will apply it in their work: to compute the homology of fibered spaces in algebraic topology [Bro59, Pro11], to construct tensor products of string field theories [GZ97], or to consider the product of Fukaya $\mathcal{A}_\infty$-categories in symplectic geometry [Sei08, Amo17]. Second, there is the community of operad theory and homotopical algebra, where the analogous result is known in the differential graded context [SU04, MS06] and expected on the topological level. Third, there are combinatorists and discrete geometers who can appreciate our result conceptually as a new development in the theory of fiber polytopes of Billera–Sturmfels [BS92].

The possible ways of iterating a binary product can be encoded, for example, by planar binary trees. Interpreting the associativity relation as an order relation, Dov Tamari introduced in his thesis [Tam51] a lattice structure on the set of planar binary trees with $n$ leaves, now known as the Tamari lattice. These lattices can be realized by polytopes, called the associahedra, in the sense that their 1-skeleton is the Hasse diagram of the Tamari lattice. We refer the reader to [CFZ02, GKPZ08, Lod04] for examples and the introduction of [CSZ15] for a comprehensive survey.

For loop spaces, composition fails to be strictly associative due to the different parametrizations, but this failure can be controlled by an infinite sequence of higher homotopies. This was made precise by James D. Stasheff in his thesis [Sta63]. He introduced a family of curvilinear polytopes, called the Stasheff polytopes, whose combinatorics coincides with the associahedra. Endowing them with a suitable operad structure, that is, an algebraic way to compose operations of various arities, allowed him to establish a now famous recognition theorem for loop spaces. In Stasheff’s theory, what is important is to have a family of contractible CW-complexes, endowed with an operad structure, whose face lattice is isomorphic to the lattice of planar trees.

Stasheff’s thesis was a profound breakthrough which opened the door to the study of homotopical algebra by means of operad-like objects. It prompted, for instance, the seminal monograph of Boardman–Vogt [BV73] on the homotopy properties of algebraic structures, and the recognition of iterated loop spaces [May72]. In this direction, Peter May introduced the little disks operads, which play a key role in many domains nowadays. In dimension 1, this gives the ‘little intervals’ operad, a finite dimensional topological operad satisfying Stasheff’s theory. Its operad structure is given by scaling a configuration of intervals in order to insert it into another interval.

Thus two communities, one working on operad and homotopy theories, the other on combinatorics and discrete geometry, seem to share a common object. Until now, however, no operad structure on any family of convex polytopal realizations of the associahedra has appeared in the literature. One was proposed in [MSS02, Part II §1.6] but does not hold as faces cannot be scaled in the same way as little intervals, and in [AA17] the problem was solved up to a notion of ‘(quasi)-normal equivalence’.
In general, the set-theoretic diagonal of a polytope will fail to be cellular. Therefore, there is a need to find a cellular approximation to the diagonal, that is, a cellular map from the polytope to its cartesian square homotopic to the diagonal. For a face-coherent family of polytopes, that is to say, a family where each face of each polytope in the family is combinatorially a product of lower-dimensional polytopes from the family, finding a family of diagonals compatible with the combinatorics of faces is a very constrained problem. In the case of the first face-coherent family of polytopes, the geometric simplices, such a diagonal map is given by the classical Alexander–Whitney map of [EZ53, EML54]. This seminal object in algebraic topology allows one to define the associative cup product on the singular cochains of a topological space. (The lack of commutativity of the cup product gives rise to the celebrated Steenrod squares [Ste47].) The next family is given by cubes, for which a coassociative approximation to the diagonal is straightforward, see Jean-Pierre Serre’s thesis [Ser51]. The associahedra form the face-coherent family of polytopes that comes next in terms of further truncations of the simplices or of combinatorial complexity. For this family there was, until now, no known approximation to the diagonal. While a face of a simplex or a cube is a simplex or a cube of lower dimension, a face of an associahedron is a product of associahedra of lower dimensions; this makes the problem of the approximation of the diagonal more intricate.

The two-fold main result of this paper is: an explicit operad structure on the Loday realizations of the associahedra together with a compatible approximation to the diagonal (Theorem 1). To accomplish this, we first consider a geometric definition (Definition 10) for a diagonal map of a polytope suitably oriented by a vector in general position. Such an approach comes from the theory of fiber polytopes of Billera–Sturmfels [BS92], after Gel’fand–Kapranov–Zelevinsky’s theory of secondary polytopes [GKZ08]. In order to define the operad structure, we resolve the issue that a face of an associahedron may not be affinely equivalent to a product of the lower-dimensional associahedra, by introducing a notion of Loday realization with arbitrary weights. Since we are looking for an operad structure compatible with the
diagonal, we define it using the diagonal, without loss of generality. (Notice that the aforementioned coherence for the diagonal maps with respect to the combinatorics of faces amounts precisely to this compatibility with the operad structure.) In the end, this provides the literature with the first object common to both of the aforementioned communities, providing discrete geometers an extra algebraic structure on realizations of the associahedra, and homotopy theorists a polytopal (and thus finite) cellular topological $A_\infty$-operad that recognizes loop spaces.

Throughout the paper, there is a dichotomy between pointwise and cellular formulas. In order to investigate their relationship and to make precise the various face-coherent properties, we introduce a meaningful notion of category of polytopes with subdivision which suits our needs. Since the definition of the diagonal maps comes from the theory of fiber polytopes, we get an induced polytopal subdivision of the associahedra. In fact, we prove a magical formula for it, in the words of Jean-Louis Loday: it is made up of the pairs of cells of matching dimensions and comparable under the Tamari order (Theorem 2). This recovers the differential graded formula of [SU04, MS06].

The new methods introduced in the present paper should allow one to attack the problem of the approximation of the diagonal for other families of polytopes, such as the ones coming from the theory of operads. Our first subsequent plan is to treat the case of the multiplihedra [Sta70] since these polytopes encode the notion of $A_\infty$-morphisms between $A_\infty$-algebras. This will provide us with a functorial construction of the tensor products of $A_\infty$-categories, which is needed in symplectic geometry. There are then the cases of the cyclohedra, permutoassociahedra, nestohedra, hypergraph polytopes, etc. These would give rise, for instance, to a tensor product construction for homotopy operads. Another relevant question the present approach allows one to study is “what kind of monoidal $\infty$-category structure does the collection of $A_\infty$-algebras admit?” In [MS06], it is proved that the differential graded diagonal cannot be coassociative. We expect that the fiber polytope method can measure the failure of this coassociativity and a useful formulation for the attacking this problem.

**Layout.** — The paper is organized as follows. The first section recalls the main relevant notions, introduces the new category of polytopes in which we work. Section 2 gives a canonical definition of the diagonal map for positively oriented polytopes and states its cellular properties. In the third section, we endow the family of Loday realizations of the associahedra with a (nonsymmetric) operad structure compatible with the diagonal maps. Section 4 states and proves the magical cellular formula for the diagonal map of the associahedra.

**Conventions.** — We use the conventions and notations of [Zie95] for convex polytopes and the ones of [LV12] for operads. We consider only convex polytopes whose vertices are the extremal points; they are equivalently defined as the intersection of finitely many half-spaces or as the convex hull of a finite set of points. We simply call them polytopes; we denote their sets of vertices by $V(P)$, their face lattices by $L(P)$, and their normal fans by $\mathcal{N}_P$. 

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1. The approximation of the diagonal of the associahedra

1.1. Planar trees, Tamari lattices, and associahedra. — We consider the set $\text{PBT}_n$ of planar binary (rooted) trees with $n$ leaves, for $n \geq 1$. We read planar binary trees according to gravity, that is from the leaves to the root. The edges of a rooted tree are of three types: the internal edges are bounded by two vertices, the leaves lie at the top and the root at the bottom.

Definition 1 (Tamari order [Tam51]). — The Tamari order is the partial order, denoted by $<$, on the set of planar binary trees generated by the following covering relation

\[ t_1 \prec t_2 \prec t_3 \prec t_4, \]

where $t_i$, for $1 \leq i \leq 4$, are planar binary trees.

For every $n \geq 1$, this forms a lattice $(\text{PBT}_n, <)$. The right-leaning leaves or internal edges are the ones of type $\downarrow$ and the left-leaning leaves or internal edges are the ones of type $\uparrow$. So two trees satisfy $s \preceq t$ if and only if one goes from $s$ to $t$ by switching pairs of successive left and right-leaning edges to corresponding pairs of successive right and left-leaning edges.

![Tamari lattice](image1)

Figure 1. The Tamari lattice $(\text{PBT}_4, <)$ with minimum at the top.
We also consider the set $\mathbf{PT}_n$ of all planar trees with $n$ leaves, for $n \geq 1$. Each of these sets forms a lattice (after adjoining a minimum) under the following partial order: a planar tree $s$ is less than a planar tree $t$, denoted $s \subset t$, if $t$ can be obtained from $s$ by a sequence of edge contractions.

**Definition 2 (Associahedra).** — For any $n \geq 2$, an $(n-2)$-dimensional associahedron is a polytope whose face lattice is isomorphic to the lattice of planar trees with $n$ leaves.

\[ \text{Figure 2. A 2-dimensional associahedron.} \]

The codimension of a face is equal to the number of internal edges of the corresponding planar tree. The 1-skeleton of an associahedron gives the Hasse diagram of the Tamari lattice.

The operation of grafting a tree $t$ at the $i$th-leaf of a tree $s$ is denoted by $s \diamond_i t$. These maps endow the collections of planar (binary) trees with a non-symmetric operad structure. We denote the corolla with $n$ leaves by $c_n$, i.e., the tree with one vertex and no internal edge. The facets of an $(n-2)$-dimensional associahedron are labeled by the two-vertex planar trees $c_{p+1} \circ c_{r+1} c_q$, for $p + q + r = n$ with $2 \leq q \leq n - 1$.

**1.2. Loday realizations of the associahedra.** — An example of realization of the associahedra can be given as follows; it is a weighted generalization of the one given by Jean-Louis Loday in [Lod04]. Notice that Loday realizations produced as convex hulls are the same polytopes as Shnider–Stasheff [SS97] produced by intersections of half-spaces, or equivalently by truncations of standard simplices.

**Definition 3 (Weighted planar binary tree).** — A *weighted planar binary tree* is a pair $(t, \omega)$ made up of a planar binary tree $t \in \mathbf{PBT}_n$ with $n$ leaves having some weight $\omega = (\omega_1, \ldots, \omega_n) \in \mathbb{Z}_{>0}^n$. We call $\omega$ the *weight* and $n$ the *arity* of the tree $t$ or the *length* of the weight $\omega$.

Let $(t, \omega)$ be a weighted planar binary tree with $n$ leaves. We order its $n-1$ vertices from left to right. At the $i$th vertex, we consider the sum $\alpha_i$ of the weights of the leaves.
supported by its left input and the sum $\beta_i$ of the weights of the leaves supported by its right input. Multiplying these two numbers, we consider the associated string which defines the following point:

$$M(t, \omega) := (\alpha_1\beta_1, \alpha_2\beta_2, \ldots, \alpha_{n-1}\beta_{n-1}) \in \mathbb{R}^{n-1}.$$ 

**Definition 4 (Loday realization).** — The Loday realization of weight $\omega$ is the polytope $K_\omega := \text{conv} \{ M(t, \omega) \mid t \in \text{PBT}_n \} \subset \mathbb{R}^{n-1}$.

The Loday realization associated to the standard weight $(1, \ldots, 1)$ is simply denoted by $K_n$. By convention, we define the polytope $K_\omega$, with weight $\omega = (\omega_1)$ of length 1, to be made up of one point labeled by the trivial tree $\cdot$. 

**Figure 3.** The Loday realization $K_5$, see Proposition 2 for the definition of the $\vec{e}_i$.

In the sequel, we will need the following key properties of these polytopes. They mainly come from [Lod04] and [For08, §6].

**Proposition 1.** — The Loday realization $K_\omega$ satisfies the following properties.

1. It is contained in the hyperplane with equation

$$\sum_{i=1}^{n-1} x_i = \sum_{1 \leq k < \ell \leq n} \omega_k \omega_\ell.$$

2. Let $p + q + r = n$ with $2 \leq q \leq n - 1$. For any $t \in \text{PBT}_n$, the point $M(t, \omega)$ is contained in the half-space defined by the inequality

$$x_{p+1} + \cdots + x_{p+q-1} \geq \sum_{p+1 \leq k < \ell \leq p+q} \omega_k \omega_\ell,$$

with equality if and only if the tree $t$ can be decomposed as $t = u \circ_{p+1} v$, where $u \in \text{PBT}_{p+1+r}$ and $v \in \text{PBT}_q$.

3. The polytope $K_\omega$ is the intersection of the hyperplane of (1) and the half-spaces of (2).

4. The face lattice $(\mathcal{L}(K_\omega), \subset)$ is equal to the lattice $(\text{PT}_n, \subset)$ of planar trees with $n$ leaves.

5. Any face of a Loday realization is isomorphic to a product of Loday realizations, via a permutation of coordinates.
Proof: — This proposition is a weighted version of the results of [Lod04], except for Point (5), which actually prompted the introduction of this more general notion.

(1) This is straightforward from the definition.
(2) This is straightforward too.
(3) Let us denote by $P$ the polytope defined by the intersection of the hyperplane of (1) and the half-spaces of (2). One can see that the points $M(t, \omega)$, for $t \in \text{PBT}_n$, are vertices of $P$, since they are defined by a system of $n - 1$ independent linear equations: the one of type (1) and $n - 2$ of type (2). In the other way round, any vertex of $P$ is characterized by a system $n - 1$ independent linear equations with the one of type (1) and $n - 2$ of type (2). We claim that any pair of equations

\[
x_{p+1} + \cdots + x_{p+q-1} = \sum_{p+1 \leq k < \ell \leq p+q} \omega_k \omega_\ell
\]

and

\[
x'_{p'+1} + \cdots + x'_{p'+q'-1} = \sum_{p'+1 \leq k < \ell \leq p'+q'} \omega_k \omega_\ell
\]

of type (2) appearing here is such that the corresponding intervals

\[ [p+1, \ldots, p+q-1] \quad \text{and} \quad [p'+1, \ldots, p'+q'-1] \]

are either nested or disjoint. If not, we are in the configuration:

\[ p+1 < p'+1 \leq p+q+1 < p'+q'+1. \]

Using these equalities and the defining inequalities of $P$, one can get

\[
x_{p'+1} + \cdots + x_{p+q-1} \leq \sum_{p+1 \leq k < \ell \leq p+q} \omega_k \omega_\ell + \sum_{p'+1 \leq k < \ell \leq p'+q'} \omega_k \omega_\ell = \sum_{p+1 \leq k < \ell \leq p+q} \omega_k \omega_\ell < \sum_{p'+1 \leq k < \ell \leq p+q} \omega_k \omega_\ell,
\]

which contradicts the definition of $P$. The solution of a system of equations as above, where the defining intervals are nested or disjoint, is a point $M(t, \omega)$, with $t \in \text{PBT}_n$.

(4) Point (2) shows that the facets of $K_\omega$ correspond bijectively to two-vertex planar trees: the facet labeled by $c_{p+1+r} \circ_{p+1} c_q$, for $p+q+r = n$ and $2 \leq q \leq n-1$, is the convex hull of the points $M(t, \omega)$ associated to planar binary trees of the form $t = u \circ_{p+1} v$, for $u \in \text{PBT}_{p+1+r}$ and $v \in \text{PBT}_q$. Any face of $K_\omega$ of codimension $k$, for $0 \leq k \leq n-2$, is defined as an intersection of $k$ facets. The above description of facets gives that the set of faces of codimension $k$ is bijectively labeled by planar trees with $k$ internal edges: the face corresponding to such a planar tree $t$ is the convex hull of points $M(s, \omega)$, for $s \subset t$. With the top dimensional face labeled by the corolla $c_n$, the statement is proved.

(5) The proof of the above point shows that it is enough to treat the case of the facets. Let $p+q+r = n$ with $2 \leq q \leq n-1$. We consider the following two weights

\[ \varpi := (\omega_1, \ldots, \omega_p, \omega_{p+1}, \cdots, \omega_{p+q}, \omega_{p+q+1}, \ldots, \omega_n) \quad \text{and} \quad \tilde{\omega} := (\omega_{p+1}, \ldots, \omega_{p+q}). \]
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The image of $K_\omega \times K_\omega \hookrightarrow K_\omega$ under the isomorphism

$$\Theta : \mathbb{R}^{p+r} \times \mathbb{R}^{q-1} \longrightarrow \mathbb{R}^{n-1}$$

$$(x_1, \ldots, x_{p+r}) \times (y_1, \ldots, y_{q-1}) \longmapsto (x_1, \ldots, x_p, y_1, \ldots, y_{q-1}, x_{p+1}, \ldots, x_{p+r})$$

is equal to the facet labeled by the planar tree $c_{p+1+r} \circ c_q$.

In other words, Point (4) shows that the polytopes $K_\omega$ are realizations of the associahedra.

1.3. The category of polytopal subdivisions. — The proposed notions of category of polytopes present in the literature only allow affine maps, which is too restrictive for our purpose: the facets of the Loday realizations of associahedra with standard weights are not affinely equivalent to the product of lower realizations with standard weights. In order to introduce a more suitable category, we begin with the following definition, which extends the classical notion of simplicial complex.

**Definition 5 (Polytopal complex).** — A polytopal complex is a finite collection $\mathcal{C}$ of polytopes of $\mathbb{R}^n$ satisfying the following conditions:

1. the empty polytope $\emptyset$ is contained in $\mathcal{C}$,
2. $P \in \mathcal{C}$ implies $\mathcal{L}(P) \subset \mathcal{C}$,
3. $P, Q \in \mathcal{C}$ implies $P \cap Q \in \mathcal{L}(P) \cap \mathcal{L}(Q)$.

Any polytope $P$ gives an example of polytopal complex $\mathcal{L}(P)$ made up of all its faces. A subcomplex of a polytopal complex $\mathcal{C}$ is a subcollection $\mathcal{D} \subset \mathcal{C}$ which forms a polytopal complex. The underlying set of a collection $\mathcal{C}$ is given by the union $|\mathcal{C}| := \bigcup_{P \in \mathcal{C}} P \subset \mathbb{R}^n$.

**Definition 6 (Polytopal subdivision).** — A polytopal subdivision of a polytope $P$ is a polytopal complex $\mathcal{C}$ whose underlying set $|\mathcal{C}|$ is equal to $P$.

The face poset $\mathcal{L}(\mathcal{C})$ of a polytopal complex is defined in the obvious way. We say that two polytopal complexes are combinatorially equivalent when their face posets are isomorphic. Now let us introduce the category we will work in.

**Definition 7 (The category Poly).** — The category Poly is made up of the following data.

- **Objects**: An object is a $d$-dimensional polytope $P$ in the $n$-dimensional Euclidian space $\mathbb{R}^n$, for any $0 \leq d \leq n$.
- **Morphisms**: A continuous map $f : P \to Q$ is a morphism when it sends $P$ homeomorphically to the underlying set $|\mathcal{D}|$ of a polytopal subcomplex $\mathcal{D} \subset \mathcal{L}(Q)$ of $Q$ such that $f^{-1}(\mathcal{D})$ defines a polytopal subdivision of $P$.

There exists obvious forgetful functors from the category Poly to the category Top of topological spaces with continuous maps and to the category CW of CW complexes with cellular maps. The latter functor is well-defined since any morphism in Poly is automatically cellular. An isomorphism $P \cong Q$ in this category is a cell-respecting homeomorphism which induces a combinatorial equivalence $\mathcal{L}(P) \cong \mathcal{L}(Q)$. 

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Lemma 1. — The category Poly endowed with the direct product × and the zero-dimensional polytope made up of one point is a symmetric monoidal category.

Proof. — The verification of the axioms is straightforward. □

This extra structure allows one to consider operads in the category Poly. Since the cellular chain functor Poly → dgMod_q is strong symmetric monoidal, it induces a functor from the category of polytopal (non-symmetric) operads to the category of differential graded (non-symmetric) operads.

1.4. The approximation of the diagonal of the associahedra. — In the sequel, we solve the following two-fold problem.

Problem
(1) Endow the collection of Loday realizations of the associahedra \( \{K_n\}_{n \geq 1} \) with a nonsymmetric operad structure in the category Poly, whose induced set-theoretical nonsymmetric operad structure on the set of faces coincides with that of planar trees.
(2) Endow the collection \( \{K_n\}_{n \geq 1} \) with diagonal maps \( \{\triangle_n : K_n \to K_n \times K_n\}_{n \geq 1} \) which form a morphism of nonsymmetric operads in the category Poly.

Remark 1
(1) Even in the category of topological spaces and for any family of realizations of the associahedra, we do not know any solution to this question in the existing literature.
(2) The compatibility of the diagonal maps with the operad structure amounts precisely to the coherence required by the approximation of the diagonal maps with respect to sub-faces by Point (5) of Proposition 1.

In order to find an operadic cellular approximation to the diagonal of the associahedra, we introduce ideas coming from the theory of fiber polytopes [BS92] as follows.

2. Canonical diagonal for positively oriented polytopes

2.1. Positively oriented polytopes

Definition 8 (Positively oriented polytope)
(1) An oriented polytope is a polytope \( P \subset \mathbb{R}^n \) endowed with a vector \( \vec{v} \in \mathbb{R}^n \) such that no edge of \( P \) is perpendicular to \( \vec{v} \), see Figure 4.

![Figure 4. An oriented polytope.](image)
(2) A positively oriented polytope is an oriented polytope \((P, \vec{v})\) such that the intersection polytope
\[ (P \cap \rho_z P, \vec{v}) \]
is oriented, where \(\rho_z := 2z - P\) stands for the reflection with respect to any point \(z \in P\), see Figure 5.

The data of an orientation vector induces a poset structure on the set of vertices \(\mathcal{V}(P)\) of \(P\), which is the transitive closure of the relation induced by the oriented edges of the 1-skeleton.

**Definition 9** (Well-oriented realization of the associahedron). — A well-oriented realization of the associahedron is a positively oriented polytope which realizes the associahedron and such that the orientation vector induces the Tamari lattice on the set of vertices.

**Proposition 2.** — Let \(\omega\) be a weight of length \(n\). Any vector \(\vec{v} = (v_1, \ldots, v_{n-1}) \in \mathbb{R}^{n-1}\) satisfying \(v_1 > v_2 > \cdots > v_{n-1}\) induces a well-oriented realization \((K_\omega, \vec{v})\) of the associahedron.

**Proof.** — Let us first prove that such an orientation vector \(\vec{v}\) induces the Tamari lattice. Let \(s < t \in \text{Tamari lattice}\). The corresponding edge in \((K_\omega, \vec{v})\) is oriented. We denote by \(\vec{z} := 2\vec{v} - \vec{n}\) the normal vectors of a facet \(F\), given by Points (1) and (2) of Proposition 1. Since edges of \(K_\omega \cap \rho_z K_\omega\) are one-dimensional intersections of facets of \(K_\omega\) or \(\rho_z K_\omega\), their directions \(\vec{d}\) are the unique solutions to a system of equations of type \((\vec{n}, \vec{d}) = 0\) and \((\vec{z}, \vec{d}) = 0\), for some set of facets \(F\). The first equation imposes \(\vec{d} = \sum_{i=1}^{n-1} a_i \vec{e}_j\), where \(\vec{e}_j := (0, \ldots, 1, -1, \ldots, 0)\), in which 1 is in the \(j\)-th place for \(1 \leq j \leq n - 2\). The second equation is equivalent to one of the following three constraints \(a_{p+1} = 0\), \(a_{p+q-1} = 0\), or \(a_{p+q-1} = a_{p+1}\). Therefore, \(\vec{d}\) is of the form \(\lambda(\vec{e}_{j_1} + \cdots + \vec{e}_{j_k})\), with \(\lambda \in \mathbb{R} \setminus \{0\}\), and so \((\vec{d}, \vec{v}) \neq 0\).

2.2. Construction of diagonal maps. — Before getting into specific argument on the associahedra, we construct a diagonal map \(\Delta : P \to P \times P\) for any positively oriented polytope \((P, \vec{v})\). Let \(\text{top}_\vec{v} P\) (resp. \(\text{bot}_\vec{v} P\)) denote the top (resp. bottom) vertex of \(P\) with respect to the orientation vector \(\vec{v}\).

**Definition 10** (Diagonal of a positively oriented polytope). — The diagonal of a positively oriented polytope \((P, \vec{v})\) is defined by
\[
\Delta_{(P, \vec{v})} : P \to P \times P \\
z \mapsto (\text{bot}_\vec{v}(P \cap \rho_z P), \text{top}_\vec{v}(P \cap \rho_z P)).
\]
Let $\beta$ be the middle-point map $P \times P \to P; (x,y) \mapsto \frac{x+y}{2}$. With the notation $\text{pr}_i$ for the projection to the $i$-th coordinate, we have $\text{pr}_1 \beta^{-1}(z) = \text{pr}_2 \beta^{-1}(z) = P \cap \rho_z P$. The diagonal $\triangle_{(P,\vec{v})}$ is a section of $\beta$ since $P \cap \rho_z P$ is symmetric with respect to the point $z$.

Since the Loday realizations $K_\omega$ of the associahedra are positively oriented when the orientation vector $\vec{v}$ has decreasing coordinates, the above formula endows them with diagonal maps, which do not depend on the choice of such an orientation vector.

**Proposition 3.** — Let $\vec{v}$ and $\vec{w}$ be two vectors of $\mathbb{R}^{n-1}$ with decreasing coordinates, then

$$\triangle_{(K_\omega,\vec{v})} = \triangle_{(K_\omega,\vec{w})},$$

for every weight $\omega$ of length $n$.

**Proof.** — The argument given in the proof of Proposition 2 shows that the formula

$$(\text{bot}_\vec{v}(K_\omega \cap \rho_z K_\omega), \text{top}_\vec{v}(K_\omega \cap \rho_z K_\omega)) = (\text{bot}_\vec{w}(K_\omega \cap \rho_z K_\omega), \text{top}_\vec{w}(K_\omega \cap \rho_z K_\omega))$$

produces the same pair for any orientation vector with decreasing coordinates. \qed

We denote by $\triangle_\omega : K_\omega \to K_\omega \times K_\omega$ the diagonal map given by any such orientation vector and we use the simple notation top and bot for associahedra. In the case of the standard weight $\omega = (1, \ldots, 1)$, we denote the diagonal simply by $\triangle_n : K_n \to K_n \times K_n$.

**Lemma 2.** — Any face $F$ of a positively oriented polytope $(P,\vec{v})$ is positively oriented once equipped with the orientation vector $\vec{v}$. Moreover, the two diagonals agree: $\triangle_{(P,\vec{v})}(z) = \triangle_{(F,\vec{v})}(z)$, for any $z \in F$.

**Proof.** — This follows from the relation $P \cap \rho_z P = F \cap \rho_z F$, for any $z \in F$. \qed
2.3. Polytopal subdivision induced by the diagonal. — The above formula for the diagonal $\triangle$ actually defines a morphism in the category Poly. To prove this property, we use some ideas coming from the theory of fiber polytopes [BS92], see also [Zie95, Chap. 9].

Let $\pi : Q \to P$ be an affine projection of polytopes with $P \subset \mathbb{R}^p$ and $Q \subset \mathbb{R}^q$. We denote by $Q_z := Q \cap \pi^{-1}(z)$ the fiber above $z \in P$. To any linear form $\psi \in (\mathbb{R}^q)^*$, we associate a collection $\mathcal{F}_\psi \subset \mathcal{L}(Q)$ as follows. We first factor the affine projection as

$$\pi = \text{pr} \circ \tilde{\pi},$$

where $\tilde{\pi} := (\pi, \psi) : Q \to \tilde{P}$, and $\text{pr} : \tilde{P} \to P; (z, t) \mapsto z$, where $\tilde{P} := \{(\pi(x), \psi(x)) \mid x \in Q\} \subset \mathbb{R}^{p+1}$.

Next let $\mathcal{L}^i(\tilde{P}) \subset \mathcal{L}(\tilde{P})$ be the subcomplex of lower faces, i.e., $F \in \mathcal{L}^i(\tilde{P})$ if and only if any $(z, t) \in F$ satisfies the equation $t = \min \psi(Q_z)$. Since the preimage of any face by a projection of polytopes is again a face, this defines a collection

$$\mathcal{F}_\psi := \{Q \cap \tilde{\pi}^{-1}(F) \mid F \in \mathcal{L}^i(\tilde{P})\} \subset \mathcal{L}(Q).$$

(This is in general not a polytopal complex since it is not stable under faces.) In the case of the point $P = \{x\}$, the unique face contained in $\mathcal{F}_\psi$ is by definition

$$Q^\psi := \{x \in Q \mid \psi(x) = \min \psi(Q)\}.$$

**Proposition 4.** — The collection $\mathcal{F}_\psi \subset \mathcal{L}(Q)$ satisfies the following properties.

1. The polytopal complex $\pi(\mathcal{F}_\psi) := \{\pi(F) \mid F \in \mathcal{F}_\psi\}$ is a polytopal subdivision of $P$.
2. For any $z \in P$, the fiber satisfies $(\mathcal{F}_\psi)_z := \pi^{-1}(z) \cap |\mathcal{F}_\psi| = (Q_z)^\psi$.

**Proof**

1. By definition $\pi(\mathcal{F}_\psi) = \pi(\mathcal{L}^i(\tilde{P}))$. The right-hand side defines a polytopal subdivision of $P$, since the restriction of the map $\text{pr}$ is a linear homeomorphism from a polytopal complex.
2. This is clear from the definition.

**Definition 11** (Coherent and tight subdivisions)

1. A subcollection $\mathcal{F} \subset \mathcal{L}(Q)$ is called a coherent subdivision of $P$ when it is of the form $\mathcal{F}_\psi$ for some $\psi \in (\mathbb{R}^q)^*$.
2. A coherent subdivision $\mathcal{F}$ is called tight when, the faces $F$ and $\pi(F)$ have the same dimension, for every $F \in \mathcal{F}$.

A coherent subdivision $\mathcal{F}$ is tight if and only if, for any $z \in P$, the fiber $\mathcal{F}_z = (Q_z)^\psi$ is a point, by Point (2) of Proposition 4. (This is also equivalent to $\mathcal{F}$ being a polytopal complex.) Therefore a tight coherent subdivision can be identified with the unique piecewise-linear section of $\pi|_Q$ which minimizes $\psi$ in each fiber. By the Point (1) of Proposition 4, this section is a morphism of the category Poly.

We apply these results to the projection

$$\beta : P \times P \to P; \quad (x, y) \mapsto \frac{x + y}{2},$$

and to the linear form $\psi(x, y) := \langle x - y, \vec{u} \rangle$ in order to obtain the following proposition.
Proposition 5. — If \((P, \vec{v})\) is a positively oriented polytope, the diagonal map \(\triangle_{(P, \vec{v})} : P \to P \times P\) is a morphism in the category Poly.

Proof. — For any \(z \in P\), the fiber \(\beta^{-1}(z)\) is the set of pairs \((x, y) \in P \times P\) such that \(x + y = 2z\). Since the sum of \(x\) and \(y\) is constant, \(\psi(x, y)\) is minimized in the fiber when \((x, \vec{v})\) is minimized, or equivalently, \((y, \vec{v})\) is maximized. On both coordinates, \(\beta^{-1}(z)\) projects down to the intersection \(P \cap (2z - P)\), which is oriented by definition. So the fiber \(\beta^{-1}(z)\) admits a unique minimal element \((\text{bot}_{\vec{v}}(P \cap \rho_z P), \text{top}_{\vec{v}}(P \cap \rho_z P))\) with respect to \(\psi\). In the end, the section defined by the tight coherent subdivision agrees with the definition of the diagonal map \(\triangle_{(P, \vec{v})}\) given in Definition 10. □

Corollary 1. — The image of \(\triangle_{(P, \vec{v})}\) is contained in \(\text{sk}_n(P \times P)\), where \(n\) is the dimension of \(P\). In particular, if one of two components of \(\triangle(z)\) lies in the interior of \(P\), then the other component is either \(\text{top}_{\vec{v}} P\) or \(\text{bot}_{\vec{v}} P\).

Remark 2. — Notice that the diagonal map \(\triangle_{(P, \vec{v})}\) is fiber-homotopic to the usual diagonal \(x \mapsto (x, x)\).

We denote by \(\mathcal{F}_{(P, \vec{v})}\) the corresponding tight coherent subdivision of \(P\) and by \(\beta(\mathcal{F}_{(P, \vec{v})})\) the polytopal subdivision of \(P\). In the case of the Loday realizations of the associahedra, Proposition 3 shows that they do not depend on the orientation vector (with decreasing coordinates), so we use the simple notations \(\mathcal{F}_{\omega}\) and \(\beta(\mathcal{F}_{\omega})\).

There is a simple way to draw this polytopal subdivision: one glues together two copies of \(P\) along an axis of direction \(\vec{v}\), then one draws all the middle points of pairs of points coming from some choices of a face on the left-hand side copy and a face on the right-hand side copy, see Figure 6. In the case of the associahedra, these choices are given by the magical formula of Section 4.

![Figure 6. Example of polytopal subdivision.](image)

Example 1. — This approach allows us to recover the classical cases of the simplices and the cubes.

(1) The classical Alexander–Whitney approximation to the diagonal of simplices \(\text{AW}_n : \Delta^n \to \Delta^n \times \Delta^n\) can be recovered with the following geometric realizations

\[
\Delta^n := \text{conv}\{(1, \ldots, 1, 0, \ldots, 0) \in \mathbb{R}^n\} = \{(z_1, \ldots, z_n) \in \mathbb{R}^n \mid 1 \geq z_1 \geq \cdots \geq z_n \geq 0\}.
\]
As usual, we denote by \( i \) the point of \( \mathbb{R}^n \) of coordinates \((1, \ldots, 1, 0, \ldots, 0)\), where 1s appear \( i \) times. If we consider the vector \( \vec{n} := (1, \ldots, 1) \), as in the proof of Proposition 2, the same argument as there shows that \((\Delta^n, \vec{n})\) is positively oriented. For \( z = (z_1, \ldots, z_n) \) satisfying \( 1 \geq z_1 \geq \cdots \geq z_i \geq 1/2 \geq z_{i+1} \geq \cdots \geq z_n \geq 0 \), one can easy see that the minimum of \( \psi \) is attained by \( \Delta(\Delta^n, \vec{n})(z) = (2z_1 - 1, \ldots, 2z_i - 1, 0, \ldots, 0, 1, \ldots, 1, 2z_{i+1}, \ldots, 2z_n) \).

We consider the faces given by the convex hull of the first \( i \) and of the last \( n - i \) vertices:

\[
\Delta^{(0, \ldots, i)} = \{ (z_1, \ldots, z_i, 0, \ldots, 0) \in \mathbb{R}^n \mid 1 \geq z_1 \geq \cdots \geq z_i \geq 0 \}
\]

and

\[
\Delta^{(i, \ldots, n)} = \{ (1, \ldots, 1, z_{i+1}, \ldots, z_n) \in \mathbb{R}^n \mid 1 \geq z_{i+1} \geq \cdots \geq z_n \geq 0 \}.
\]

The tight coherent subdivision of \( \Delta^n \) is equal to

\[
\mathcal{F}(\Delta^n, \vec{n}) = \{ \Delta^{(0, \ldots, i)} \times \Delta^{(i, \ldots, n)} \mid 0 \leq i \leq n \},
\]

which recovers the (simplicial) Alexander–Whitney map of [EZ53, EML54].

(2) The approximation of the diagonal \( C^n \to C^n \times C^n \) of the cube \( C^n : = [0, 1]^n \) used by Jean-Pierre Serre in [Ser51] can easily be recovered by the present method. First, it is a positively oriented polytope once equipped with the orientation vector \( \vec{n} \).

Since an \( n \)-dimensional cube is a product of \( n \) intervals, the various formulas are straightforward.

Notice that these two examples work particularly well because we do not need to stretch the faces and their combinatorial complexity is very limited: any face appearing here is affinely equivalent to a lower dimensional polytope of the respective family. These properties do not hold anymore for the Loday realizations of the associahedra; such a difficulty is omnipresent in the rest of this paper.

3. Operad structure on Loday realizations

3.1. Cellular properties

**Proposition 6.** — Suppose \( P \) and \( Q \) are normally equivalent polytopes, i.e., with same normal fans \( \mathcal{N}_P = \mathcal{N}_Q \). If \( P \) and \( Q \) are well-oriented by the same orientation vector \( \vec{v} \), then the tight coherent subdivisions \( \mathcal{F}(P, \vec{v}) \) and \( \mathcal{F}(Q, \vec{v}) \) are combinatorially equivalent in a canonical way.

**Proof.** — Normal equivalence of polytopes induces combinatorial equivalence

\[
\Phi : \mathcal{L}(P) \cong \mathcal{L}(Q).
\]

By the definition of the diagonal map, a pair of points \((x, y) \in P \times P\) is contained in the image \( \text{Im} \Delta_P \) if and only if it satisfies the following condition: there exists no vector \( \vec{w} \) and positive number \( \varepsilon > 0 \) with \( \langle \vec{v}, \vec{w} \rangle > 0 \) and \( x - \varepsilon \vec{w}, y + \varepsilon \vec{w} \in P \). The
latter conditions can be restated in terms of normal cones as follows. Recall that, for any subset $C \subset \mathbb{R}^n$, the polar cone $C^*$ of $C$ is defined by

$$C^* := \{x \in \mathbb{R}^n \mid \forall x \in C, \langle x, y \rangle \leq 0\}. $$

The polar cone theorem asserts that $C^{**}$ is the smallest closed convex cone which contains $C$. By definition, $(P - y)^*$ is the normal cone $\mathcal{N}_P(G)$ corresponding to the face which satisfies $y \in \hat{G}$. Applying the polar cone theorem to $C = P - y$, we obtain that $(P - y)^{**}$ is the normal cone $\mathcal{N}_P(G)$ corresponding to the face which satisfies $y \in \hat{G}$. Applying the polar cone theorem to $C = P - y$, we obtain that $(P - y)^{**}$ is the set of vectors $\vec{w}$ such that $y + \varepsilon \vec{w} \in P$ for some $\varepsilon > 0$. Therefore if $x \in \hat{F}$ and $y \in \hat{G}$, the condition for $(x, y) \in \text{Im } \Delta_P$ is that there exists no vector $\vec{w}$ such that $\langle \vec{v}, \vec{w} \rangle > 0$ and $\vec{w} \in -\mathcal{N}_P(F)^* \cap \mathcal{N}_P(G)^*$. Since this condition depends only on the normal fan, the map $\Phi \times \Phi : \mathcal{L}(P \times P) \to \mathcal{L}(Q \times Q)$ induces the canonical combinatorial equivalence $\mathcal{F}(P, \vec{v}) \cong \mathcal{F}(Q, \vec{w})$.

**Corollary 2.** Let $\omega$ and $\theta$ be two weights of same length. The two polytopal subdivisions $\beta(\mathcal{F}_\omega)$ and $\beta(\mathcal{F}_\theta)$ of $K_\omega$ and $K_\theta$ respectively are combinatorially equivalent, i.e., labeled by the same pairs of planar trees.

**Proof.** This is a direct corollary of Point (3) of Proposition 1 and Proposition 6.

Proposition 3 and Corollary 2 show that the type of faces composing the polytopal subdivision of the Loday realizations of the associahedra are intrinsic: they depend neither on the orientation vector (with decreasing coordinates) nor on the weight. From now on, we simply denote them by $\mathcal{F}_n \subset \mathcal{F} \times \mathcal{F}$ and $\beta(\mathcal{F}_n)$. Their description will be the subject of the magical formula given in Section 4.

**3.2. Pointwise properties.** We can enhance the above one-to-one correspondence of polytopal subdivisions to the pointwise level using the isomorphism in the category Poly.

**Proposition 7.** Let $(P, \vec{v})$ and $(Q, \vec{w})$ be two positively oriented polytopes, with a combinatorial equivalence $\Phi : \mathcal{L}(P) \cong \mathcal{L}(Q)$. Suppose that tight coherent subdivisions $\mathcal{F}(P, \vec{v})$ and $\mathcal{F}(Q, \vec{w})$ are combinatorially equivalent under $\Phi \times \Phi$.

1. There exists a unique continuous map $\text{tr} = \text{tr}_{P}^{Q} : P \longrightarrow Q$, which extends the restriction $\mathcal{V}(P) \to \mathcal{V}(Q)$ of $\Phi$ to the set of vertices and which commutes with the respective diagonal maps.

2. The map $\text{tr}$ is an isomorphism in the category Poly, whose correspondence of faces agrees with $\Phi$.

**Proof**

1. In the core of this proof, we use the simple notation $\Delta_P$ for $\Delta_{(P, \vec{v})}$ and

$$\Delta_P^{(n)} := \Delta_P^{2^n-1} \circ \Delta_P^{2^n-2} \circ \cdots \circ \Delta_P^2 \circ \Delta_P : P \longrightarrow P^{2^n}.$$
for its iterations. We also denote the averaging map by
\[ \beta_P^{(n)} : P^{2^n} \longrightarrow P \]
\[ (x_1, \ldots, x_{2^n}) \mapsto \frac{x_1 + \cdots + x_{2^n}}{2^n}. \]

Notice that \( \Delta_P^{(n)} \) is a section of \( \beta_P^{(n)} \). Any map \( \text{tr} \) commuting with the respective diagonal maps satisfies
\[ \text{tr} = \beta_Q^{(n)} \circ \text{tr}^{2^n} \circ \Delta_P^{(n)}. \]

Let us prove the statement by induction on the dimension \( d \) of the polytopes \( P \) and \( Q \). It is obvious for \( d = 0 \). For \( d = 1 \), let us suppose that \( P = Q = [0,1] \) and that \( \Phi(0) = 0, \Phi(1) = 1 \), without any loss of generality. Since the two polytopal subdivisions correspond bijectively under \( \Phi \), the definition of the diagonal maps shows that \( \bar{v} \) and \( \bar{w} \) are oriented in the same direction. By dichotomy, one can check that the \( 2^n \)-tuple \( \Delta_P^{(n)}(k/2^n) \) is made up of \( 2^n - k \) zeros and \( k \) ones. This induces the formula
\[ \text{tr}(k/2^n) = \beta_Q^{(n)} \circ \Phi^{2^n} \circ \Delta_P^{(n)}(k/2^n) = k/2^n. \]

The identity of \([0,1]\) therefore provides the map \( \text{tr} \), and is unique by continuity.

Let us now suppose that the statement holds up to dimension \( d - 1 \) and let \( P \) and \( Q \) be two polytopes of dimension \( d \). Since the restriction of the diagonal map of \( P \) to a face \( F \in \mathcal{L}(P) \) is equal to the diagonal map of the face, i.e., \( \Delta_{(F,\sigma)}(z) = \Delta_{(F,\sigma)}(z) \), for any \( z \in F \), by Lemma 2, the induction hypothesis implies that the transition map \( \text{tr} \) exists and is uniquely defined on the \((d-1)\)-skeleton of \( P \). To study the interior of the top face of \( P \), we consider the following filtration
\[ P(n) := P \supset \bigcup_{k=0}^{2^n-1} k \text{bot } P + (2^n - 1 - k) \text{top } P + P, \]
for \( n > 0 \). Notice that
\[ \bigcup_{n>0} P(n) = P \supset \text{bot } P, \text{top } P, \]
where \([\text{bot } P, \text{top } P]\) stands for the line segment defined by the top and the bottom vertices of \( P \). Corollary 1 shows that the faces appearing in the tight coherent subdivision corresponding to the section \( \Delta_P^{(n)} \) are of two kinds:

* (\text{bot } P, \ldots, \text{bot } P, \text{top } P, \ldots, \text{top } P)
* or \((F_1, \ldots, F_{2^n})\), with codim \( F_i \geq 1 \), for all \( i \) such that \( 1 \leq i \leq 2^n \).

The images of the first ones under \( \beta_P^{(n)} \) are equal to the sets excluded from \( P \) in the definition of \( P(n) \). Otherwise stated, the image of any \( z \in P(n) \) under the iterated diagonal map satisfies \( \Delta_P^{(n)}(z) \in (\text{sk}_{d-1} P)^{2^n} \), and so
\[ \text{tr}(z) = \beta_Q^{(n)} \circ (\text{tr}|_{\partial P})^{2^n} \circ \Delta_P^{(n)}(z). \]

The image of the transition map \( \text{tr} \) on the main axis \([\text{bot } P, \text{top } P]\) is given by the same dichotomy argument as in the case \( d = 1 \). In the end, this proves the uniqueness of the transition map.

To show the existence of such a suitable transition map, we define it by the above-mentioned formulas. It remains to prove the continuity at points \( x \in [\text{bot } P, \text{top } P] \).
of the main axis. Let $\varepsilon > 0$ and let us find $\delta > 0$ such that $|x - y| < \delta$ implies $|\text{tr}(x) - \text{tr}(y)| < \varepsilon$. We consider $n \in \mathbb{Z}_{>0}$ satisfying $\text{diam} \, Q < 2^n \varepsilon$. There are two cases to consider.

(a) When $x$ cannot be written as $(k \text{ bot } P + (2^n - k) \text{ top } P)/2^n$, with $0 \leq k \leq 2^n$, it is of the form $(k \text{ bot } P + (2^n - 1 - k) \text{ top } P + x')/2^n$ for some $x' \in P$. In this case, we can take a small enough $\delta > 0$ such that $|x - y| < \delta$ implies that $y$ is of the form $(k \text{ bot } P + (2^n - 1 - k) \text{ top } P + y')/2^n$, for some $y' \in P$. This implies

- $\Delta_P^{(n)}(x) = (\text{ bot } P, \ldots, \text{ bot } P, x', \text{ top } P, \ldots, \text{ top } P)$
- and $\Delta_P^{(n)}(y) = (\text{ bot } P, \ldots, \text{ bot } P, y', \text{ top } P, \ldots, \text{ top } P)$,

with the same number of bot $P$ and top $P$, and so

$$|\text{tr}(x) - \text{tr}(y)| = \frac{|\text{tr}(x') - \text{tr}(y')|}{2^n} \leq \frac{\text{diam} \, Q}{2^n} < \varepsilon.$$

(b) Otherwise, we can write $x$ as $(k \text{ bot } P + (2^n - k) \text{ top } P)/2^n$, with $0 \leq k \leq 2^n$.

We further divide into the following two cases and take the least $\delta$.

(i) When $y \not\in P(n)$, we can take $\delta > 0$ such that if $|x - y| < \delta$, then $y$ is contained in

- $(k \text{ bot } P + (2^n - 1 - k) \text{ top } P + P)/2^n$
- or $((k - 1) \text{ bot } P + (2^n - k) \text{ top } P + P)/2^n$.

This implies $|\text{tr}(x) - \text{tr}(y)| < \varepsilon$ as above.

(ii) When $y \in P(n)$, observe that $\beta_Q^{(n)} \circ (\text{tr}|_{\partial P})^{2^n} \circ \Delta_P^{(n)}$ actually defines a continuous restriction of tr to the closed set

$$P \setminus \left( \bigcup_{k=0}^{2^n-1} k \text{ bot } P + (2^n - 1 - k) \text{ top } P + P \right) /2^n,$$

which contains both $x$ and $y$. Therefore we can choose $\delta$ which satisfies the condition.

(2) This is straightforward from the above description. The inverse morphism in $\text{Poly}$ is the transition map $\text{tr}_Q^P$. \hfill $\square$

Corollary 3. — Any two normally equivalent polytopes positively oriented by the same orientation vector $\vec{v}$ are isomorphic in the category $\text{Poly}$, with an isomorphism commuting with the diagonal maps.

Proof. — This is a direct corollary of Proposition 6 and Proposition 7. \hfill $\square$

This produces a stronger comparison between the diagonal maps of two Loday realizations associated to different weights than Corollary 2: the transition map $\text{tr} = \text{tr}_Q^P : K_{\omega} \to K_{\theta}$ preserves homeomorphically the faces of the same type and it commutes with the respective diagonals. Up to isomorphisms in the category $\text{Poly}$, the diagonal maps do not depend on the orientation vector (with decreasing coordinates) nor on the weight.
3.3. Operad structure. — We use the above results to endow the collection \( \{K_n\}_{n \geq 1} \) of Loday realizations (with standard weights) with an operad structure as follows.

Definition 12 (Operad structure). — For any \( n, m \geq 1 \) and any \( 1 \leq i \leq m \), we define the partial composition map by

\[
\circ_i : K_m \times K_n \xrightarrow{\text{tr} \times \text{id}} K_{(1, \ldots, n, \ldots, 1)} \times K_n \xrightarrow{\Theta} K_{n+m-1},
\]

where the last inclusion is given by the block permutation of the coordinates introduced in the proof of Proposition 1.

Theorem 1

(1) The collection \( \{K_n\}_{n \geq 1} \) together with the partial composition maps \( \circ_i \) form a non-symmetric operad in the category Poly.

(2) The maps \( \{\triangle_n : K_n \to K_n \times K_n\}_{n \geq 1} \) form a morphism of non-symmetric operads in the category Poly.

Proof. — By Proposition 7 and Proposition 5, the various maps are morphisms in the category Poly.

(1) We need to prove the sequential and the parallel composition axioms of a non-symmetric operad, see [LV12, §5.3.4]. The sequential composition axiom amounts to the commutativity of the following diagram

\[
\begin{array}{c}
K_\ell \times K_m \times K_n \xrightarrow{\text{id} \times \circ_j} K_\ell \times K_{m+n-1} \\
\downarrow \circ_i \times \text{id} \quad \quad \downarrow \circ_i \\
K_{\ell+m-1} \times K_n \xrightarrow{\circ_i \circ_j - 1} K_{\ell+m+n-2}.
\end{array}
\]

Let us denote by \( F \) the face of \( K_{\ell+m+n-2} \) labeled by the composite tree

\[
\circ_\ell \circ_i (c_m \circ_j c_n) = (c_\ell \circ_i c_m) \circ_{i+j-1} c_n.
\]

The two maps of this diagram have the same image equal to \( F \). They both induce two cellular homeomorphisms \( K_\ell \times K_m \times K_n \to F \) which meet the requirements of Proposition 7 by Proposition 1 and by the fact that the transition map \( \text{tr} \) and the isomorphism \( \Theta \) commute with the diagonal maps. So they are equal by Point (1) of Proposition 7. The parallel composition axiom is proved in the same way.

(2) This statement means that the partial composition maps commute with the diagonal maps, which is the case since the maps \( \text{tr} \) and \( \Theta \) do.

Under the cellular chain functor, we recover the classical differential graded non-symmetric operad \( \mathcal{A}_\infty \) encoding homotopy associative algebras [Sta63], see also [LV12, Chap. 9]. To understand the induced diagonal on the differential graded level, we need the following magical formula describing its cellular structure.
4. The magical formula

Theorem 2 (Magical formula). — For any Loday realization of the associahedra, the approximation of the diagonal satisfies

\[
\text{Im} \, \Delta_n = \bigcup_{\dim F + \dim G = n-2} F \times G.
\]

Figure 7. The polytopal subdivision \( \beta (\mathcal{F}_4) \) of \( K_4 \).

The pairs of faces appearing on the right-hand side of the magical formula are called matching pairs. In other words, the tight coherent subdivision

\[
\mathcal{F}_n = \{(F,G) \mid \text{top } F \leq \text{bot } G, \dim F + \dim G = n-2\},
\]

made up of matching pairs, gives a polytopal subdivision of the associahedron under \( \beta \).

Applying the cellular chain functor, we recover the differential graded diagonal given in [SU04, MS06]. Notice that neither the pointwise version nor the cellular version of the diagonal map \( \Delta_n \) can be coassociative by [MS06, Th. 13].

4.1. First step: \( \text{Im} \, \Delta_n \subset \bigcup F \times G \). — We prove this property more generally for any product \( P := K_{\omega_1} \times \cdots \times K_{\omega_k} \) of Loday realizations of the associahedra. Recall that \( P \subset \mathbb{R}^N \), where \( N := n_1 + \cdots + n_k - k \), and \( d := \dim P = n_1 + \cdots + n_k - 2k \). The
set $\mathcal{V}(P)$ of vertices of $P$ coincides with $\text{PBT}_{n_1} \times \cdots \times \text{PBT}_{n_k}$. By Proposition 2, any vector $\vec{v} = (v_1, v_2, \ldots, v_N)$ satisfying

$$v_1 > \cdots > v_{n_1-1}, \quad v_{n_1} > \cdots > v_{n_1+n_2-2}, \ldots, \quad v_{n_1+\cdots+n_k-k+2} > \cdots > v_N$$

makes $(P, \vec{v})$ into a positively oriented polytope with 1-skeleton isomorphic to the product of Tamari lattices.

We consider the map $L = (L_i)_{1 \leq i \leq m-2} : \text{PBT}_m \to \{0,1\}^{m-2}$ to the Boolean lattice defined by

$$L_i(t) := \begin{cases} 0 & \text{if the } (i+1)\text{-th leaf is left-leaning } \vee, \\ 1 & \text{if the } (i+1)\text{-th leaf is right-leaning } \wedge. \end{cases}$$

We extend it to a map $L = (L_i)_{1 \leq i \leq d} : \mathcal{V}(P) \to \{0,1\}^d$. We consider the collection of vectors $\vec{e}_1, \ldots, \vec{e}_d$ in $\mathbb{R}^N$ defined by $\vec{e}_i := (0, \ldots, 1, -1, \ldots, 0)$, where 1 is in the place $a_i$, with $\{a_i\}_{1 \leq i \leq d}$ the increasing sequence found by starting with $1, \ldots, N$ and deleting the values $n_1 + \cdots + n_j - j$, for $j = 1, \ldots, k$.

**Lemma 3.** The following properties are satisfied:

1. $s \leq t \Rightarrow L(s) \leq L(t)$,
2. $\langle \vec{e}_i, \vec{v} \rangle > 0$, for $1 \leq i \leq d$,
3. each edge connecting $s \in L_i^{-1}(0)$ and $t \in L_i^{-1}(1)$ is parallel to $\vec{e}_i$,
4. any fiber $L_i^{-1}(b)$, for $b \in \{0,1\}^d$, is contained in a facet of $P$.

**Proof.**

1. When we switch a pair of successive left and right-leaning edges to a pair of successive right and left-leaning edges, either it does not change the orientation of the leaves or it just changes one left-leaning leaf into a right-leaning leaf.

2. This is straightforward from the definition of $\vec{v}$.

3. It is enough to prove it on one polytope $K_\omega$. In this case, $t$ is a planar binary tree obtained from a planar binary tree $s$ by switching the $(i+1)^{th}$ leaf from left-leaning to right-leaning, which implies

$$M(s, \omega)M(t, \omega) = \omega \vec{v}_{i+2} \vec{e}_i.$$

4. Reading the sequence $(1, b_1, \ldots, b_{n_1-2}, 0)$ from left to right, there is at least one occurrence of $(1, 0)$, say at place $i$ and $i+1$. Every face labeled by a forest of trees $t$ satisfying $L(t) = b$ lies in the facet labeled by $(c_{n_1-1} \circ c_2, c_{n_2}, \ldots, c_{n_k})$. □

**Lemma 4.** Let $F$ be a face of $P$ which contains a vertex $s$ such that $L_i(s) = 1$. For any $x \in F$, there exists $\varepsilon > 0$ such that $x - \varepsilon \vec{v}_i \in P$.

**Proof.** It is enough to perform the proof for one Loday realization $P = K_\omega$ that we endow with the linear form $\psi(x) := \langle x, \vec{e}_i \rangle$. We claim that, for the projection $id : P \to P$, the associated subcomplexes of lower and upper faces, introduced in Section 2.3, are given by

$$\mathcal{V}(\mathcal{L}^+(\tilde{P})) = \{M(t, \omega) \mid L_i(t) = 0\} \quad \text{and} \quad \mathcal{V}(\mathcal{L}^-(\tilde{P})) = \{M(t, \omega) \mid L_i(t) = 1\}.$$
The intersection $L_i^{-1}(0) \cap \mathscr{L}\hat{\mathcal{P}}(\hat{P})$ is nonempty, since it contains the vertex bot $P$.
Suppose that the polytopal complex $\mathscr{L}\hat{\mathcal{P}}(\hat{P})$ contains a vertex living in $L_i^{-1}(1)$. By the connectivity of $\mathscr{L}\hat{\mathcal{P}}(\hat{P})$, it admits an edge with vertices $s \in L_i^{-1}(0)$ and $t \in L_i^{-1}(1)$.

Point (3) of Lemma 3 says that such an edge is parallel to $\vec{e}_i$, which contradicts the minimality of $\mathscr{L}\hat{\mathcal{P}}(\hat{P})$. This proves the inclusion $\mathcal{V}(\mathscr{L}\hat{\mathcal{P}}(\hat{P})) \subset \{M(t, \omega) \mid L_i(t) = 0\}$.

The opposite inclusion is proved by the same argument as for $\mathscr{L}\hat{\mathcal{P}}(\hat{P})$ together with the fact that the union $\mathscr{L}\hat{\mathcal{P}}(\hat{P}) \cup \mathcal{L}^\uparrow(\hat{P})$ contains all the vertices of $P$. If not, a vertex $x$, which is not contained in it, is neither maximal nor minimal with respect to $\vec{e}_i$, which contradicts the fact that $x$ is an extremal point. This characterization of the polytopal complex of lower faces shows that any face $F$ of $P$ containing a vertex $s$ such that $L_i(s) = 1$ satisfies $F \not\subset \mathcal{V}(\mathscr{L}\hat{\mathcal{P}}(\hat{P}))$, and thus $F \cap |\mathscr{L}\hat{\mathcal{P}}(\hat{P})| = \emptyset$, which concludes the proof.

**Proposition 8.** — Let $F$ and $G$ be two faces of $P$ of matching dimensions, i.e., $\dim F + \dim G = d$. We consider $s := \text{top} F$ and $t := \text{bot} G$. When $L(s) \not\subseteq L(t)$, we have $F \times G \cap \text{Im } \Delta = \emptyset$.

**Proof.** — When $L(s) \not\subseteq L(t)$, there exists $i$ such that $L_i(s) = 1$ and $L_i(t) = 0$.

By Lemma 4 and a dual version of it proved with the same arguments, for every $x \in F$ and $y \in G$, there exists $\varepsilon > 0$ such that $(x - \varepsilon \vec{e}_i, y + \varepsilon \vec{e}_i) \in P \times P$. Suppose now that there exists $(x, y) \in F \times G \cap \text{Im } \Delta$. Since $(x + y)/2 = ((x - \varepsilon \vec{e}_i) + (y + \varepsilon \vec{e}_i))/2$, the two points $(x, y)$ and $(x - \varepsilon \vec{e}_i, y + \varepsilon \vec{e}_i)$ lie in the same fiber of $\beta$. As we saw in the proof of Proposition 5, the point $(x, y)$ minimizes $\langle x, \vec{v} \rangle$ in the fiber of $\beta$. However, the computation

$$\langle \vec{v}, x - \varepsilon \vec{e}_i \rangle = \langle \vec{v}, x \rangle - \varepsilon \langle \vec{v}, \vec{e}_i \rangle < \langle \vec{v}, x \rangle$$

violates the definition of $\Delta$.

Proposition 8 excludes faces $F \times G$ of matching dimensions with $L(s) \not\subseteq L(t)$ from the image of $\Delta$. Notice that, by Point (1) of Lemma 3, $L(s) \not\subseteq L(t)$ implies $s \not\subseteq t$.

In order to exclude the remaining case when $s \not\subseteq t$ and $L(s) \subseteq L(t)$, we prepare the following two lemmas.

**Lemma 5.** — Let $t \in \text{PBT}_{n_1} \times \cdots \times \text{PBT}_{n_k}$ be a forest of planar binary trees with a total of $r_t + k$ right-leaning leaves and $\ell_t + k$ left-leaning leaves. There exists a unique maximal (with respect to inclusion) face $F_t$ (resp. $G_t$) of $P$ such that $\text{top } F = t$ (resp. $\text{bot } G = t$). The dimensions of these faces are given by $\dim F_t = \ell_t$ and $\dim G_t = r_t$.

**Proof.** — The cell $F_t$ (resp. $G_t$) can be obtained by collapsing all the left-leaning (resp. right-leaning) internal edges of all the trees of the forest $t$. One can see that, for any forest of planar binary trees, the number of left-leaning (resp. right-leaning) internal edges is equal to the number of right-leaning (resp. left-leaning) leaves minus $k$. 

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**Lemma 6.** — Let \( F, G \subset P \) be a pair of faces of matching dimensions. When \( s := \text{top } F \) and \( t := \text{bot } G \) satisfy \( L(s) \leq L(t) \), then \( (F, G) = (F_s, G_t) \) and \( L(s) = L(t) \).

**Proof.** — By definition, \( F \subset F_s \) and \( G \subset G_t \), and thus \( \dim F + \dim G \leq \ell_s + r_t \leq d \) by Lemma 5. The top dimension assumption \( \dim F + \dim G = d \) allows us to conclude that \( \dim F = \ell_s, \dim G = r_t, F = F_s, G = G_t \), and \( L(s) = L(t) \). □

We can now conclude Step 1: prove by induction on the dimension \( d \) of \( P \) that any pair \( F, G \subset P \) of faces of matching dimensions with \( s \not< t \) and \( L(s) \leq L(t) \) satisfies \( F \times G \not\subset \text{Im } \triangle \). This is straightforward to check in dimensions \( d = 0 \) and \( d = 1 \). In dimension \( d \), let us suppose, to the contrary, that there exists such a pair \( F, G \) of faces satisfying \( F \times G \subset \text{Im } \triangle \). Lemma 6 implies \( L(s) = L(t) \). In this case, both \( s \) and \( t \) lie in a common facet \( Q \) of \( P \). By Point (4) of Lemma 3. By Lemma 2, the induced diagonal \( \triangle_{(Q,\vec{v})} \) on \( Q \) is the restriction of \( \triangle_{(P,\vec{v})} \). Let us consider \( F' := F \cap Q \) and \( G' := G \cap Q \). If \( \dim F' + \dim G' > \dim Q \), then \( F' \times G' \not\subset \triangle_{(Q,\vec{v})} \). By Corollary 1. When \( \dim F' + \dim G' \leq \dim Q \), we consider any pair of faces \( F' \subset F'' \) and \( G' \subset G'' \) of \( Q \) of matching dimensions: \( \dim F'' + \dim G'' = \dim Q \). They satisfy \( r := \text{top } F'' \geq \text{top } F' = s \) and \( u := \text{bot } G'' \leq \text{bot } G' = t \), which implies \( r \not< u \).

We claim that \( F'' \times G'' \not\subset \text{Im } \triangle_{(Q,\vec{v})} \). Since \( Q \) is a facet of \( P \), it is of type \( Q' := K_{\omega_1} \times K_{\omega_2} \times \cdots \times K_{\omega_k} \cong Q \), under the isomorphism \( \Theta : \mathbb{R}^N \cong \mathbb{R}^N \) given by the block permutation of coordinates described in the proof of Point (5) of Proposition 1. The image of the orientation vector \( \vec{v} = (v_1, \ldots, v_{n_1-1}, v_{n_1}, \ldots) \) under the inverse permutation of coordinates is equal to

\[
\vec{v}' := \Theta^{-1}(\vec{v}) = (v_1, \ldots, v_{i-1}, v_{i+m-1}, \ldots, v_{n_1-1}, v_i, \ldots, v_{i+m-2}, v_{n_1}, \ldots),
\]

so it well-orient \( Q' \). Therefore the isomorphism \( \Theta \) intertwines the two diagonals \( \triangle_{(Q,\vec{v})} \) and \( \triangle_{(Q',\vec{v}')} \). Let us denote by \( r = (r_1, \ldots, r_k) \) and \( u = (u_1, \ldots, u_k) \) the various planar trees composing the two forests. Under the notation of the proof of Point (5) of Proposition 1, one can write the two planar trees \( r_1 = \vec{r}_1 \circ_1 \vec{r}_1 \) and \( u_1 = \vec{u}_1 \circ_1 \vec{u}_1 \). This gives

\[
r' := \text{top } \Theta^{-1}(F'') = (\vec{r}_1, \vec{r}_2, \ldots, r_k)
\]

and

\[
u' := \text{bot } \Theta^{-1}(G'') = (\vec{u}_1, \vec{u}_2, \ldots, u_k).
\]

We have \( r' \not< u' \). Indeed, the condition \( r \not< u \) implies that there exists \( 1 \leq j \leq k \) such that \( r_j \not< u_j \) and \( r_j \not< u_j \). If \( j = 1 \), then automatically \( r' \not< u' \). If \( j = 1 \), then either \( \vec{r}_1 \not< \vec{u}_1 \) or \( \vec{r}_1 \not< \vec{u}_1 \), since \( \circ_1 \) preserves the Tamari order, and again \( r' \not< u' \). If \( L(r') \not< L(u') \), we conclude with Proposition 8, otherwise we conclude our claim with the induction hypothesis. Finally, notice that any cell appearing in \( \text{Im } \triangle \) is contained in a product of cells of matching dimensions. The above argument shows that \( F' \times G' \) cannot appear in \( \text{Im } \triangle_{(Q,\vec{v})} \). This concludes the first step.

4.2. **Second step:** \( \text{Im } \triangle_n \supset \bigcup F \times G. \) — In this section, we prove that every matching pair \( (F, G) \) of \( K_n \) satisfies \( F \times G \subset \text{Im } \triangle_n \). By Point (1) of Lemma 3 and by Lemma 6 such pairs are of type \( (F_s, G_t) \) with \( s \leq t \) and \( L(s) = L(t) \).
Proposition 9. — For every \( t \in \text{PBT}_n \), we have \( F_t \times G_t \subset \text{Im} \triangle_n \).

Proof. — By the pointwise formula \((t,t) \in \text{Im} \triangle_n\). But \( F_t \times G_t \) is the only cell coming from a matching pair that can contain \((t,t)\) by Section 4.1. \( \square \)

Lemma 7. — Let \( t \prec u \) be an edge of \( K_n \) satisfying \( L(t) = L(u) \).

1. The cell \( G_t \cap G_u \) is a facet of \( G_t \) and \( G_u \).
2. The cell \( G_t \cap G_u \) is not of type \( G_v \), for \( v \in \text{PBT}_n \).
3. If \( G_t \cap G_u \) is a facet of \( G_v \), then \( v = t \) or \( v = u \).

Proof

(1) Let the planar binary tree \( u \) be obtained from \( t \) by switching a pair \((f,e)\) of successive left and right-leaning edges to a pair \((f',e')\) of successive right and left-leaning edges:

\[
t = f \quad f' \quad t_2 \quad t_3 \\
\quad t_4
\]

By Lemma 5, the cell \( G_t \cap G_u \) corresponds to a tree \( s \) obtained by collapsing all the right-leaning internal edges of \( u \) except \( f' \). The cell \( G_u \) corresponds to the tree \( s/f' \) obtained from \( s \) by contracting the edge \( f' \). The cell \( G_t \) corresponds to the tree \( s/e' \).

(2) Notice first that a cell \( G \) is of type \( G_v \), for \( v \in \text{PBT}_n \) if and only if it is labeled by a planar tree having only left-leaning internal edges. This is not the case of the tree \( s \), since it carries the right-leaning internal edge \( f' \).

(3) In this case, the labeling tree \( r \) of \( G_v \) is obtained from \( s \) by contracting one internal edge. Since \( r \) should not contain any right-leaning internal edge, it can only be obtained in two ways: by contracting \( e' \) or \( f' \). In the former case \( v = t \) and in the latter case \( v = u \). \( \square \)

Proposition 10. — If \( s \leq t \) and \( L(s) = L(t) \) then \( F_s \times G_t \subset \text{Im} \triangle_n \).

Proof. — We fix \( s \) and we consider the sub-lattice \( \text{PBT}^{>s}_n \) of elements greater than \( s \). Let us prove by induction on the elements \( t \) of \( \text{PBT}^{>s}_n \cap L^{-1}(L(s)) \) that \( F_s \times G_t \subset \text{Im} \triangle_n \). The base case is \( s = t \) and this is done in the above Proposition 9. Suppose the conclusion holds for \( t \) and let \( t \prec u \) be an edge satisfying \( L(t) = L(u) \). Then \( H := F_s \times G_t \cap F_s \times G_u = F_s \times (G_t \cap G_u) \) is a facet of \( F_s \times G_t \) by Point (1) of Lemma 7 and it does not project into the boundary of \( K_n \) via \( \beta \). Since \( F_s \times G_t \) lies in \( \text{Im} \triangle_n \), it is attached along \( H \) to a cell \( F_r \times G_v \) by Section 4.1. We claim that \( r = s \), \( v = u \). Since \( H = F_s \times (G_t \cap G_u) \) is a facet of \( F_r \times G_v \), there are two options: either \( F_s = F_r \) or \( G_t \cap G_u = G_v \). The latter case is contradicted by Point (2) of Lemma 7. Therefore, \( G_t \cap G_u \) is a facet of \( G_v \), and we conclude by Point (3) of Lemma 7. \( \square \)

This concludes the proof of Step 2.
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Naruki Masuda, Johns Hopkins University, Department of Mathematics
3400 N. Charles Street, Baltimore, MD 21218, USA
E-mail: nmasuda2@jhu.edu
Url: https://math.jhu.edu/~nmasuda2/

Hugh Thomas, Département de mathématiques, Université du Québec à Montréal
Local PK-5151, 201, Avenue du Président-Kennedy, Montréal, Canada
E-mail: thomas.hugh_r@uqam.ca
Url: https://math.uqam.ca/corps-professoral/professeurs/professeurs-reguliers/professeur/thomas.hugh_r/

Andy Tonks, Department of Mathematics, University of Leicester,
University Road, Leicester LE1 7RH, United Kingdom
E-mail: apt12@leicester.ac.uk
Url: https://www2.le.ac.uk/departments/mathematics/extranet/staff-material/staff-profiles/andy-tonks

Bruno Vallette, Laboratoire Analyse, Géométrie et Applications, Université Sorbonne Paris Nord,
CNRS, UMR 7539
93430 Villetaneuse, France
E-mail: vallette@math.univ-paris13.fr
Url: https://www.math.univ-paris13.fr/~vallette/